

Two-dimensional Kripke Semantics

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Formal Structures for Computation and Deduction (FSCD 2024)

Talinn, 13 July 2024



Kripke semantics vs. type theory

Modal logic is important in Computer Science:

- ▶ temporal logic
- ▶ epistemic logic
- ▶ dynamic logic
- ▶ Hennessy-Milner logic

In most cases, it is given a **Kripke semantics**.

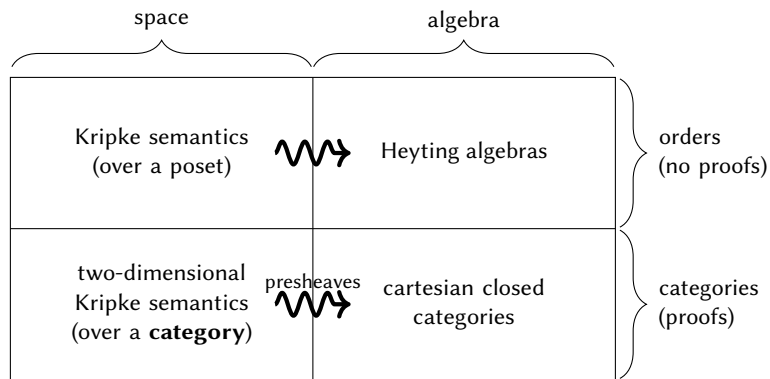
But in type theory **proofs are important** (Curry-Howard-Lambek).

Type-theoretic **modalities** arise *everywhere*:

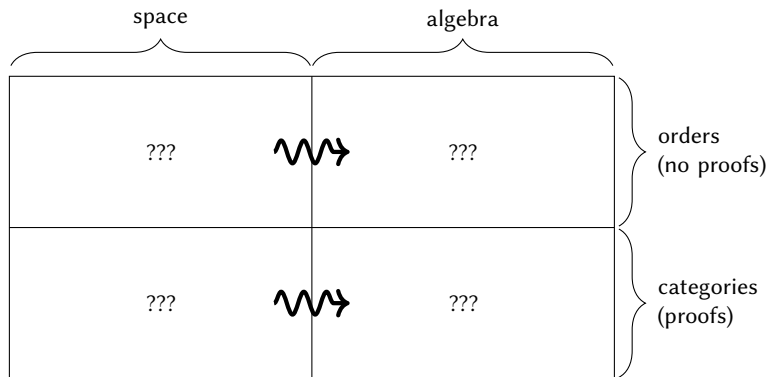
- ▶ ‘logical’ time
- ▶ proof-irrelevance
- ▶ globality
- ▶ information flow

How can we connect these two worlds?

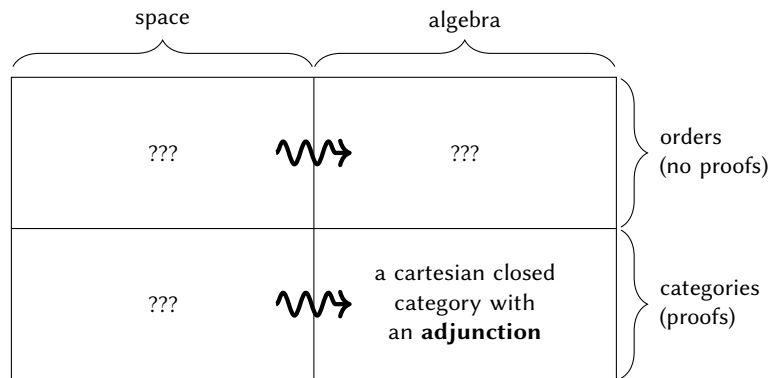
Models of intuitionistic logic



Models of intuitionistic modal logic



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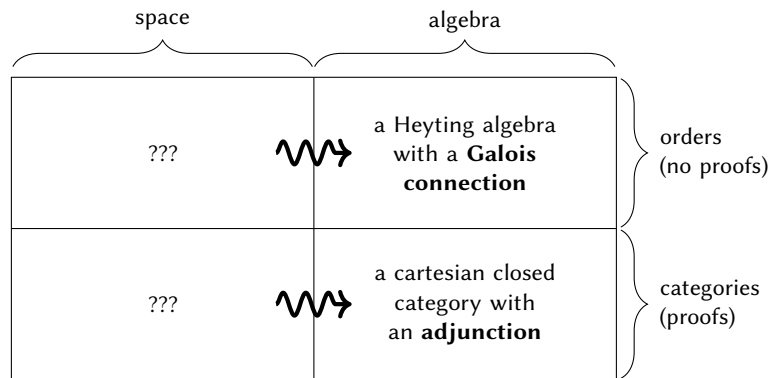


Using an adjunction was proposed by Clouston [Clo18].

It has proven remarkably robust in modal type theory.

This is an **objective** answer in the sense of Lawvere [Law94; LS09].

Models of intuitionistic modal logic

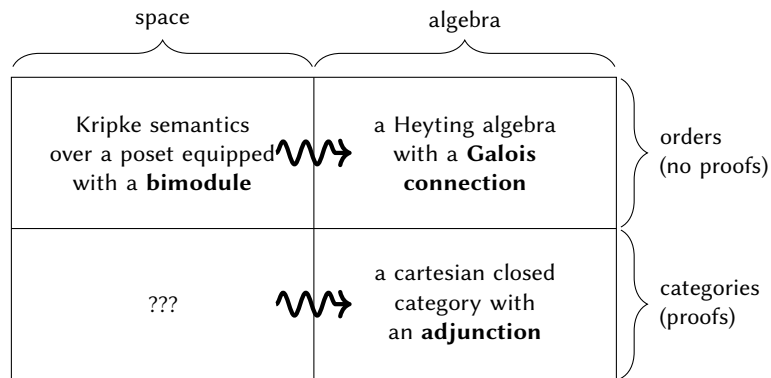


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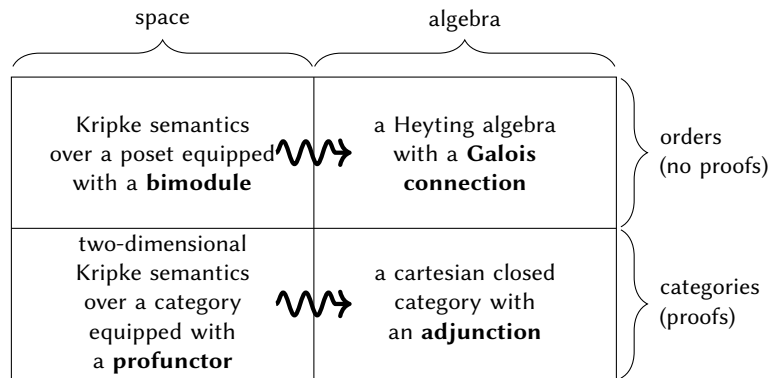


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Kripke semantics of intuitionistic logic

Let (W, \sqsubseteq) be a **Kripke frame**, i.e. a partial order.

$\text{Up}(W) =$ **upper sets** $S \subseteq W$ (where $w \in S$ and $w \sqsubseteq v$ imply $v \in S$)

Let $V : \text{Var} \rightarrow \text{Up}(W)$ map each proposition to an upper set. Define

$$w \vDash p \stackrel{\text{def}}{\equiv} w \in V(p)$$

$$w \vDash \perp \stackrel{\text{def}}{\equiv} \text{never}$$

$$w \vDash \varphi \wedge \psi \stackrel{\text{def}}{\equiv} w \vDash \varphi \text{ and } w \vDash \psi$$

$$w \vDash \varphi \vee \psi \stackrel{\text{def}}{\equiv} w \vDash \varphi \text{ or } w \vDash \psi$$

$$w \vDash \varphi \rightarrow \psi \stackrel{\text{def}}{\equiv} \forall v. w \sqsubseteq v \text{ and } v \vDash \varphi \text{ imply } v \vDash \psi$$

Monotonicity: $w \vDash \varphi$ and $w \sqsubseteq v$ imply $v \vDash \varphi$

Theorem (Kripke)

A formula is valid (in all frames and all worlds) iff it is a theorem.

Algebraic semantics of intuitionistic logic

A **Heyting algebra** (H, \leq) is a lattice (has finite meets and joins) such that for every $x, y \in H$ there exists $x \Rightarrow y \in H$ with

$$c \wedge x \leq y \iff c \leq x \Rightarrow y \quad \text{for all } c \in H$$

Suppose that for each proposition p we have an element $\llbracket p \rrbracket \in H$. Intuitionistic logic can then be interpreted into H compositionally:

$$\begin{aligned}\llbracket \perp \rrbracket &\stackrel{\text{def}}{=} 0 \\ \llbracket \varphi \wedge \psi \rrbracket &\stackrel{\text{def}}{=} \llbracket \varphi \rrbracket \wedge \llbracket \psi \rrbracket \\ \llbracket \varphi \vee \psi \rrbracket &\stackrel{\text{def}}{=} \llbracket \varphi \rrbracket \vee \llbracket \psi \rrbracket \\ \llbracket \varphi \rightarrow \psi \rrbracket &\stackrel{\text{def}}{=} \llbracket \varphi \rrbracket \Rightarrow \llbracket \psi \rrbracket\end{aligned}$$

Theorem

A formula is valid (= 1 in all algebras) iff it is a theorem.

Prime algebraic lattices: from space to algebra

Let (W, \sqsubseteq) be a **Kripke frame**, and $\mathbb{2} \stackrel{\text{def}}{=} \{0 \sqsubseteq 1\}$.

$[W, \mathbb{2}]$ (= monotone maps $W \rightarrow \mathbb{2}$) has many curious properties:

- ▶ $[W, \mathbb{2}] \cong \text{Up}(W)$ where the order is inclusion
- ▶ It is a **complete Heyting algebra** (arbitrary joins and meets)
- ▶ The **principal upper set** embedding $\uparrow : W^{\text{op}} \rightarrow [W, \mathbb{2}]$ given by $w \mapsto \{v \mid w \sqsubseteq v\}$ preserves meets and exponentials.
- ▶ An element is a **prime** ($p \sqsubseteq \bigsqcup_i d_i \Rightarrow \exists i. p \sqsubseteq d_i$) iff it is $\uparrow w$.
- ▶ Every upper set S is a join of primes:

$$S = \bigsqcup \{P \mid P \text{ prime}, P \subseteq S\} = \bigsqcup \{\uparrow w \mid w \in S\}$$

In short: $[W, \mathbb{2}]$ is a **prime algebraic lattice** [Win09].

There is a **duality**: $\text{Pos}^{\text{op}} \simeq \text{PrAlgLatt}$.

Intuitionistic logic: from space to category

Play the same trick as before, but replace $\mathcal{2}$ by **Set** [Law73].

The category $[\mathcal{C}, \mathbf{Set}]$ of presheaves $\mathcal{C} \rightarrow \mathbf{Set}$:

- ▶ is a **(co)complete cartesian closed category**
- ▶ The **Yoneda embedding** $\mathbf{y} : \mathcal{C}^{\text{op}} \rightarrow [\mathcal{C}, \mathbf{Set}]$ given by $\mathbf{y}(w) \stackrel{\text{def}}{=} \text{Hom}(w, -)$ preserves products and exponentials.
- ▶ A presheaf P is **tiny** just if $\text{Hom}(P, -)$ preserves colimits. All representables are tiny [and vice versa if \mathcal{C} is Cauchy-complete].
- ▶ Every presheaf $P : \mathcal{C} \rightarrow \mathbf{Set}$ is a colimit of tiny objects:

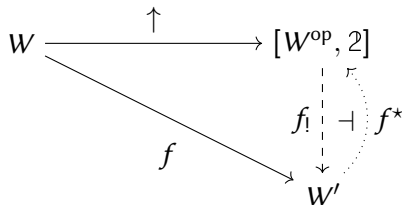
$$P = \lim_{\rightarrow (w,x) \in \text{el } P} \mathbf{y}(w)$$

There is a duality: $\mathbf{Cat}_{\text{cc}}^{\text{op}} \simeq \mathbf{PshCat}$ (Bunge's theorem).

2D Kripke semantics = semantics in $[\mathcal{C}, \mathbf{Set}]$.

Extensions

Let W' be a **complete lattice**, and let $f : W \rightarrow W'$ be monotone.



$f!$: the **unique join-preserving** map satisfying $f!(\uparrow w) = f(w)$.

$$f!(S) \stackrel{\text{def}}{=} \bigsqcup \{f(w) \mid w \in S\}$$

As both lattices are complete, this has a right adjoint f^* . Explicitly:

$$f^*(w') \stackrel{\text{def}}{=} \{w \mid f(w) \sqsubseteq w'\}$$

Then

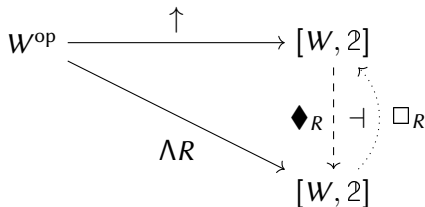
$$f!(S) \sqsubseteq w' \iff S \subseteq f^*(w')$$

Bimodules and Extensions

Let (W, \sqsubseteq) be a Kripke frame. $R \subseteq W \times W$ is a **bimodule** just if

$$w' \sqsubseteq w R v \sqsubseteq v' \implies w' R v'$$

Equivalently: $R : W^{\text{op}} \times W \rightarrow \mathcal{2}$. Now extend $\Lambda R : W^{\text{op}} \rightarrow [W, \mathcal{2}]$:



Concretely:
$$\begin{cases} \blacklozenge_R(S) \stackrel{\text{def}}{=} \{w \in W \mid \exists v. v R w \text{ and } v \in S\} \\ \square_R(S) \stackrel{\text{def}}{=} \{w \in W \mid \forall v. w R v \text{ implies } v \in S\} \end{cases}$$

Every such adjunction on $[W, \mathcal{2}]$ corresponds to a bimodule!

Duality: $\mathbf{EBimod}^{\text{op}} \simeq \mathbf{PrAlgLattO}$.

The logic of Dzik, Järvinen, and Kondo [DJK10]

A very simple **tense logic** with two modalities, \blacklozenge and \square .

Kripke semantics:

$$w \vDash \blacklozenge \varphi \stackrel{\text{def}}{\equiv} \exists v. v R w \text{ and } v \vDash \varphi$$

$$w \vDash \square \varphi \stackrel{\text{def}}{\equiv} \forall v. w R v \text{ implies } v \vDash \varphi$$

Algebraic semantics: a Heyting algebra with a Galois connection.

$$\frac{\blacklozenge \varphi \rightarrow \psi}{\varphi \rightarrow \square \psi}$$

and

$$\frac{\varphi \rightarrow \square \psi}{\blacklozenge \varphi \rightarrow \psi}$$

Some derivable rules:

$$\frac{\varphi \rightarrow \psi}{\square \varphi \rightarrow \square \psi}$$

$$\frac{\varphi}{\square \top}$$

$$\frac{}{\square \top}$$

$$\frac{\blacklozenge \perp}{\perp}$$

$$\frac{\varphi \rightarrow \psi}{\blacklozenge \varphi \rightarrow \blacklozenge \psi}$$

The usual \blacklozenge is **not monotonic** in this setting.

Lifting to categories

- ▶ Replace bimodules by **profunctors**
- ▶ Use **left Kan extension** along Yoneda

This leads to a duality $\mathbf{EProf}_{\mathcal{C}\mathcal{C}}^{\text{op}} \simeq \mathbf{PshCatO}$.

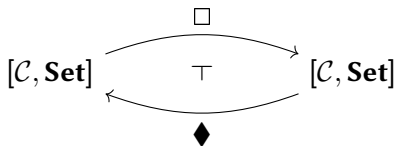
Modalities on presheaves $P : \mathcal{C} \rightarrow \mathbf{Set}$:

$$(\blacklozenge P)(w) = \int^{v \in \mathcal{C}} R(v, w) \times P(v)$$

$$(\blacksquare P)(w) = \int_{v \in \mathcal{C}} R(w, v) \rightarrow P(v) \cong \text{Hom}_{[\mathcal{C}, \mathbf{Set}]}(R(w, -), \llbracket \varphi \rrbracket)$$

Theorem

A two-dimensional Kripke semantics over \mathcal{C} uniquely corresponds to



Key facts

- ▶ The Kripke semantics (over a poset) and the algebraic semantics of intuitionistic (modal) logic are related by a **duality**.
- ▶ The interpretation of modalities arises canonically by extension (from a bimodule).
- ▶ These dualities can be **restricted** to ‘truth-preserving’ and ‘deduction-preserving’ morphisms respectively.
- ▶ They can also be **categorified**, relating 2D Kripke semantics (over a category) with cartesian closed categories.
- ▶ The interpretation of modalities arises canonically by **Kan extension** (from a profunctor).

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