

Two-dimensional Kripke Semantics

Alex Kavvos

University of Bristol

British Colloquium for Theoretical Computer Science

University of Bath

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Kripke semantics vs. type theory

Modal logic is important in Computer Science:

- ▶ temporal logic
- ▶ epistemic logic
- ▶ dynamic logic
- ▶ Hennessy-Milner logic

In most cases, it is given a **Kripke semantics**.

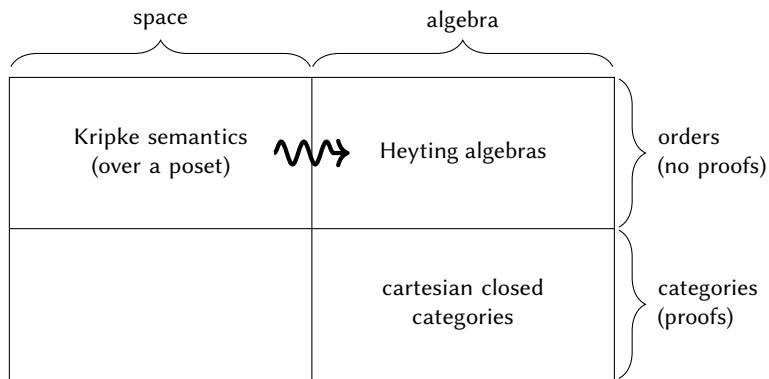
But in type theory **proofs are important** (Curry-Howard-Lambek).

Type-theoretic **modalities** arise *everywhere*:

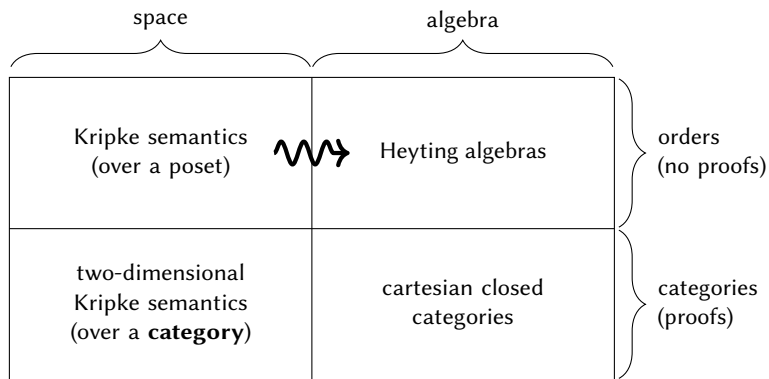
- ▶ 'logical' time
- ▶ proof-irrelevance
- ▶ globality
- ▶ information flow

How can we connect these two worlds?

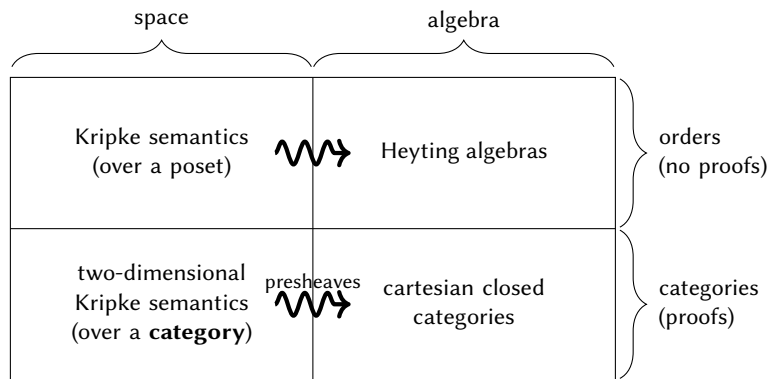
Models of intuitionistic logic



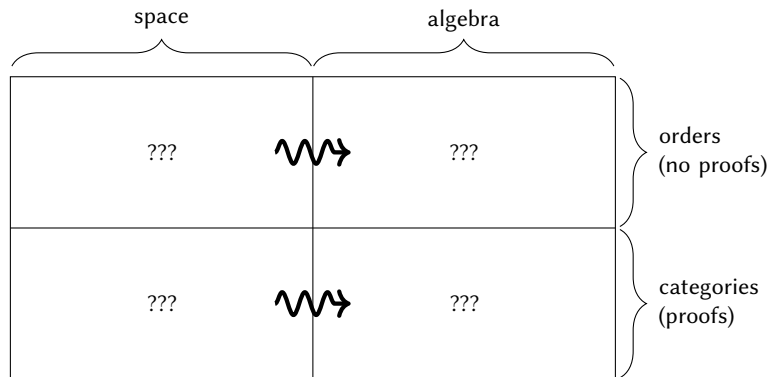
Models of intuitionistic logic



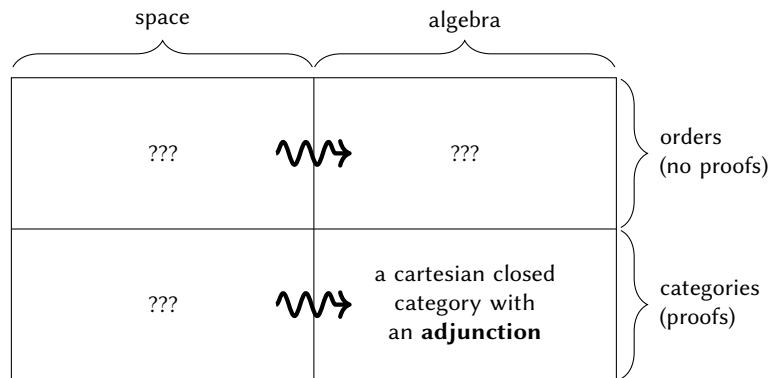
Models of intuitionistic logic



Models of intuitionistic modal logic

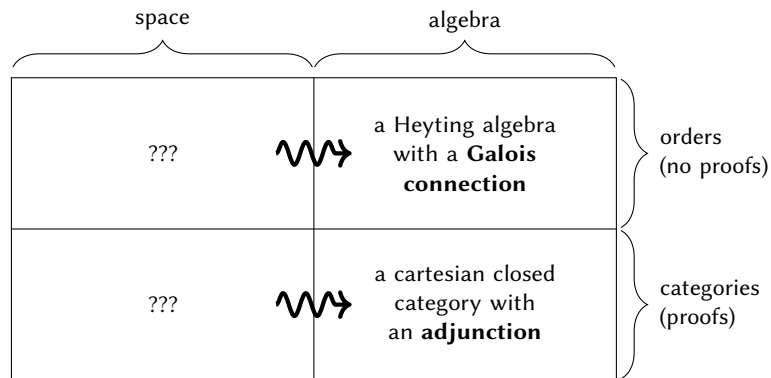


Models of intuitionistic modal logic



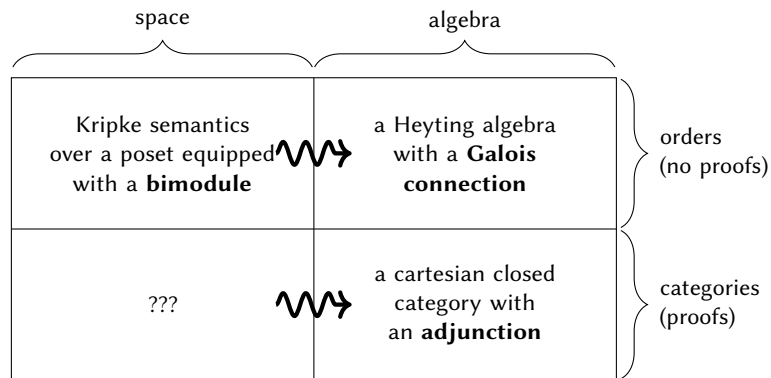
Using an adjunction was proposed by Clouston [Clo18]. It has proven remarkably robust in modal type theory. This is an **objective** answer in the sense of Lawvere [Law94; LS09].

Models of intuitionistic modal logic



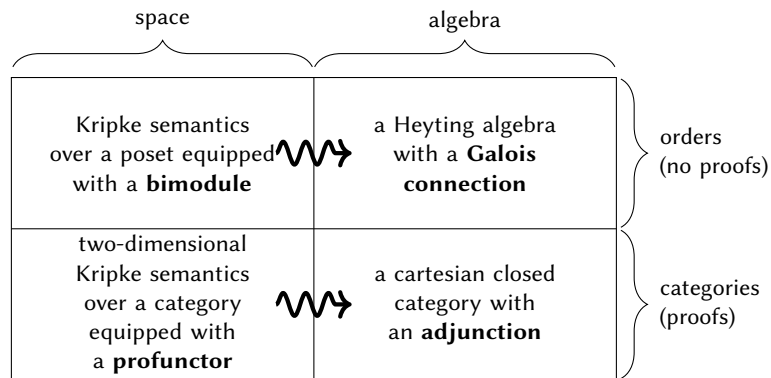
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Roadmap

Intuitionistic logic: Space vs. Algebra

Modal logic: bimodules

Coherent semantics

I. INTUITIONISTIC LOGIC: SPACE VS. ALGEBRA

Kripke semantics of intuitionistic logic

Let (W, \sqsubseteq) be a **Kripke frame**, i.e. a partial order.

$\text{Up}(W) =$ **upper sets** $S \subseteq W$ (where $w \in S$ and $w \sqsubseteq v$ imply $v \in S$)

Let $V : \text{Var} \rightarrow \text{Up}(W)$ map each proposition to an upper set. Define

$$w \vDash p \stackrel{\text{def}}{\equiv} w \in V(p)$$

$$w \vDash \perp \stackrel{\text{def}}{\equiv} \text{never}$$

$$w \vDash \varphi \wedge \psi \stackrel{\text{def}}{\equiv} w \vDash \varphi \text{ and } w \vDash \psi$$

$$w \vDash \varphi \vee \psi \stackrel{\text{def}}{\equiv} w \vDash \varphi \text{ or } w \vDash \psi$$

$$w \vDash \varphi \rightarrow \psi \stackrel{\text{def}}{\equiv} \forall v. w \sqsubseteq v \text{ and } v \vDash \varphi \text{ imply } v \vDash \psi$$

Monotonicity: $w \vDash \varphi$ and $w \sqsubseteq v$ imply $v \vDash \varphi$

Theorem (Kripke)

A formula is valid (in all frames and all worlds) iff it is a theorem.

Algebraic semantics of intuitionistic logic

A **Heyting algebra** (H, \leq) is a lattice (has finite meets and joins) such that for every $x, y \in H$ there exists $x \Rightarrow y \in H$ with

$$c \wedge x \leq y \iff c \leq x \Rightarrow y \quad \text{for all } c \in H$$

Suppose that for each proposition p we have an element $\llbracket p \rrbracket \in H$. Intuitionistic logic can then be interpreted into H compositionally:

$$\begin{aligned}\llbracket \perp \rrbracket &\stackrel{\text{def}}{=} 0 \\ \llbracket \varphi \wedge \psi \rrbracket &\stackrel{\text{def}}{=} \llbracket \varphi \rrbracket \wedge \llbracket \psi \rrbracket \\ \llbracket \varphi \vee \psi \rrbracket &\stackrel{\text{def}}{=} \llbracket \varphi \rrbracket \vee \llbracket \psi \rrbracket \\ \llbracket \varphi \rightarrow \psi \rrbracket &\stackrel{\text{def}}{=} \llbracket \varphi \rrbracket \Rightarrow \llbracket \psi \rrbracket\end{aligned}$$

Theorem

A formula is valid (= 1 in all algebras) iff it is a theorem.

From space to algebra I

Let (W, \sqsubseteq) be a Kripke frame, and $\mathcal{2} \stackrel{\text{def}}{=} \{0 \sqsubseteq 1\}$.

$[W, \mathcal{2}]$ (= monotone maps $W \rightarrow \mathcal{2}$) has many curious properties:

- ▶ $[W, \mathcal{2}] \cong \text{Up}(W)$ where the order is inclusion
- ▶ It is a **complete Heyting algebra** (arbitrary joins and meets)
- ▶ The **principal upper set** $\uparrow w \stackrel{\text{def}}{=} \{v \mid w \sqsubseteq v\}$ gives a map

$$\uparrow : W^{\text{op}} \rightarrow [W, \mathcal{2}] \quad \text{monotone, order-embedding}$$

- ▶ An element is a **prime** ($p \sqsubseteq \bigsqcup_i d_i \Rightarrow \exists i. p \sqsubseteq d_i$) iff it is $\uparrow w$.
- ▶ Every upper set S is a join of primes:

$$S = \bigsqcup \{P \mid P \text{ prime, } P \subseteq S\} = \bigsqcup \{\uparrow w \mid w \in S\}$$

In short: $[W, \mathcal{2}]$ is a **prime algebraic lattice** [Win09].

From space to algebra II

Theorem (Raney; Nielsen, Plotkin, Winskel)

Every prime algebraic lattice is isomorphic to $[W, 2]$ for some W .

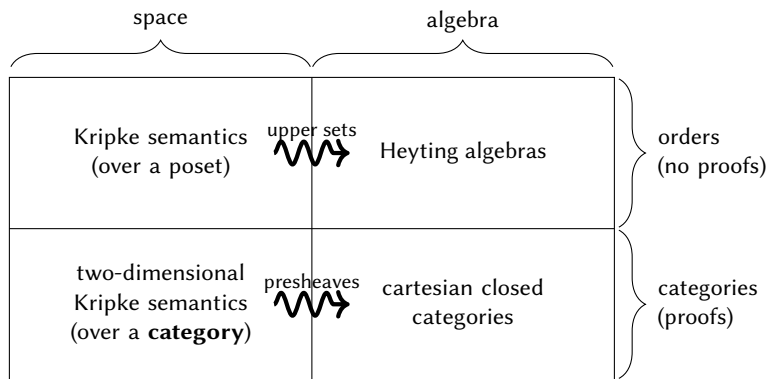
This amounts to a **duality**

$$\mathbf{Pos}^{\text{op}} \simeq \mathbf{PrAlgLatt}$$

The Kripke semantics over (W, \sqsubseteq) can be related to the HA $[W, 2]$:

$$w \models \varphi \iff w \in \llbracket \varphi \rrbracket$$

Models of intuitionistic logic



Categories as spaces

A category \mathcal{C} has

- ▶ **objects** $c, d, \dots \in \mathcal{C}$
- ▶ **morphisms** $f, g, \dots : c \rightarrow d$ between two objects
- ▶ a way to **compose** morphisms, and **identity** morphisms

Categories are often used as ‘mathematical universes’
(sets, graphs, vector spaces, topological spaces, ...)

But also, a category can be seen as a **partial order with evidence**.

$$\frac{}{x \sqsubseteq x}$$

$$\frac{x \sqsubseteq y \quad y \sqsubseteq z}{x \sqsubseteq z}$$

$$\frac{}{\text{id}_x : x \rightarrow x}$$

$$\frac{f : x \rightarrow y \quad g : y \rightarrow z}{g \circ f : x \rightarrow z}$$

A category can also be seen as a *space with direction*.

Two-dimensional Kripke semantics of intuitionistic logic

Take any (small) category \mathcal{C} . Given arbitrary sets $\llbracket p \rrbracket_w$, define a set

$$\llbracket \varphi \rrbracket_w$$

of **proofs** of φ at a world $w \in \mathcal{C}$, by induction on φ .

$$\llbracket \perp \rrbracket_w \stackrel{\text{def}}{=} \emptyset$$

$$\llbracket \varphi \wedge \psi \rrbracket_w \stackrel{\text{def}}{=} \llbracket \varphi \rrbracket_w \times \llbracket \psi \rrbracket_w = \{(x, y) \mid x \in \llbracket \varphi \rrbracket_w, y \in \llbracket \psi \rrbracket_w\}$$

$$\llbracket \varphi \vee \psi \rrbracket_w \stackrel{\text{def}}{=} \llbracket \varphi \rrbracket_w + \llbracket \psi \rrbracket_w = \{(1, a) \mid a \in \llbracket \varphi \rrbracket_w\} \cup \{(2, b) \mid b \in \llbracket \psi \rrbracket_w\}$$

$$\llbracket \varphi \rightarrow \psi \rrbracket_w \stackrel{\text{def}}{=} (v : \mathcal{C}) \rightarrow \text{Hom}_{\mathcal{C}}(w, v) \rightarrow \llbracket \varphi \rrbracket_v \rightarrow \llbracket \psi \rrbracket_v \quad (\text{not exactly})$$

Monotonicity: for each $f : w \rightarrow v$ and $x \in \llbracket \varphi \rrbracket_w$ define $f \cdot x \in \llbracket \varphi \rrbracket_v$

This defines a **presheaf**

$$\llbracket \varphi \rrbracket : \mathcal{C} \longrightarrow \mathbf{Set}$$

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From space to category

Play the same trick as before, but replace $\mathbb{2}$ by **Set** [Law73].

The category $[\mathcal{C}, \mathbf{Set}]$ of presheaves $\mathcal{C} \longrightarrow \mathbf{Set}$:

- ▶ is a **(co)complete cartesian closed category**
- ▶ **Representables** $\mathbf{y}(w) \stackrel{\text{def}}{=} \text{Hom}(w, -)$ give an embedding

$$\mathbf{y} : \mathcal{C}^{\text{op}} \longrightarrow [\mathcal{C}, \mathbf{Set}]$$

- ▶ A presheaf P is **tiny** just if $\text{Hom}(P, -)$ preserves colimits. All representables are tiny (and vice versa if \mathcal{C} is Cauchy-complete).
- ▶ Every presheaf $P : \mathcal{C} \longrightarrow \mathbf{Set}$ is a colimit of tiny objects:

$$P = \lim_{\longrightarrow (w,x) \in \text{el } P} \mathbf{y}(w)$$

From space to algebra I

Let (W, \sqsubseteq) be a Kripke frame, and $\mathcal{2} \stackrel{\text{def}}{=} \{0 \sqsubseteq 1\}$.

$[W, \mathcal{2}]$ (= monotone maps $W \rightarrow \mathcal{2}$) has many curious properties:

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In short: $[W, \mathcal{2}]$ is a **prime algebraic lattice** [Win09].

Duality

This construction gives us a duality

$$\mathbf{Cat}_{\text{cc}}^{\text{op}} \simeq \mathbf{PshCat}$$

between

- ▶ (Cauchy-complete, small) categories (\approx ‘2D’ Kripke frames)
- ▶ presheaf categories (\approx proof-relevant prime alg. lattices)

In short:

A two-dimensional Kripke semantics is a categorical semantics in a presheaf category $[\mathcal{C}, \mathbf{Set}]$.

II. MODAL LOGIC: BIMODULES

What is intuitionistic modal logic?

Many possible answers, in particular around $\diamond!$ [DM23]

The Simpson [Sim94] criteria:

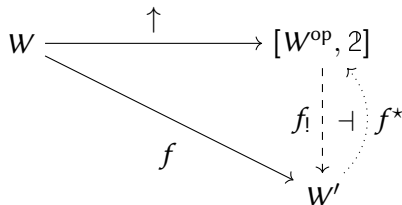
1. It should be conservative over intuitionistic logic.
2. It should prove all intuitionistic theorems (even with modalities).
3. Adding $\varphi \vee \neg\varphi$ should yield a classical modal logic.
4. It should satisfy the disjunction property.
5. \Box and \diamond should be independent.
6. Its semantics should be ‘intuitionistically comprehensible.’

#6 is formalised by translation to intuitionistic first-order logic.

An alternative proposal: **let category theory show you the way.**

Extensions

Let W' be a **complete lattice**, and let $f : W \rightarrow W'$ be monotone.



$f_!$: the **unique join-preserving** map satisfying $f_!(\uparrow w) = f(w)$.

$$f_!(S) \stackrel{\text{def}}{=} \bigsqcup \{f(w) \mid w \in S\}$$

As both lattices are complete, this has a right adjoint f^* . Explicitly:

$$f^*(w') \stackrel{\text{def}}{=} \{w \mid f(w) \sqsubseteq w'\}$$

Then

$$f_!(S) \sqsubseteq w' \iff S \subseteq f^*(w')$$

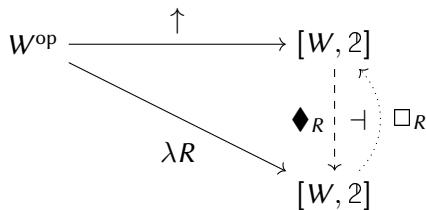
Bimodules and Extensions

Let (W, \sqsubseteq) be a Kripke frame. How to add an accessibility relation?

$R \subseteq W \times W$ is a **bimodule** just if

$$w' \sqsubseteq w R v \sqsubseteq v' \implies w' R v'$$

Equivalently: $R : W^{\text{op}} \times W \rightarrow \mathcal{2}$. Now extend $\Lambda R : W^{\text{op}} \rightarrow [W, \mathcal{2}]$:



Concretely:

$$\blacklozenge_R(S) \stackrel{\text{def}}{=} \{w \in W \mid \exists v. v R w \text{ and } v \in S\}$$

$$\square_R(S) \stackrel{\text{def}}{=} \{w \in W \mid \forall v. w R v \text{ implies } v \in S\}$$

The logic of Dzik, Jarvinen, and Kondo [DJK10]

A very simple **tense logic** with two modalities, \blacklozenge and \Box .

Kripke semantics:

$$w \vDash \blacklozenge \varphi \stackrel{\text{def}}{\equiv} \exists v. v R w \text{ and } v \vDash \varphi$$

$$w \vDash \Box \varphi \stackrel{\text{def}}{\equiv} \forall v. w R v \text{ implies } v \vDash \varphi$$

Algebraic semantics: a Heyting algebra with a Galois connection.

$$\frac{\blacklozenge \varphi \rightarrow \psi}{\varphi \rightarrow \Box \psi}$$

and

$$\frac{\varphi \rightarrow \Box \psi}{\blacklozenge \varphi \rightarrow \psi}$$

Some derivable rules:

$$\frac{\varphi \rightarrow \psi}{\Box \varphi \rightarrow \Box \psi}$$

$$\frac{\varphi}{\Box \top}$$

$$\frac{}{\Box \top}$$

$$\frac{\blacklozenge \perp}{\perp}$$

$$\frac{\varphi \rightarrow \psi}{\blacklozenge \varphi \rightarrow \blacklozenge \psi}$$

The usual \blacklozenge is **not monotonic** in this setting.

Lifting to categories

As before we get a duality: $\mathbf{EBimod}^{\text{op}} \simeq \mathbf{PrAlgLattO}$.

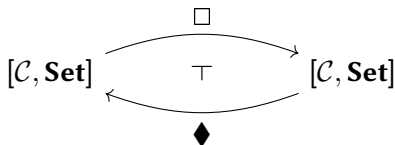
To categorify:

- ▶ Replace bimodules by **profunctors**
- ▶ Use **left Kan extension** along Yoneda

This leads to a duality $\mathbf{EProf}_{\text{cc}}^{\text{op}} \simeq \mathbf{PshCatO}$.

Theorem

A two-dimensional Kripke semantics of this logic corresponds to



But adjunctions on a cartesian closed category were exactly the models of the modal λ -calculus of Clouston [Clo18]!

III. COHERENT SEMANTICS

Completeness?

The developments so far only prove **relative completeness**:

- ▶ Suppose a formula is valid in all Heyting algebras.
- ▶ Then it is valid in all prime algebraic lattices.
- ▶ Then it is valid in all Kripke semantics

∴ the algebraic semantics is as complete as the Kripke semantics.

How to get the opposite direction?

The classic proof (Gehrke and van Gool [Gv24, §4.4]):

- ▶ Make a Kripke frame of **prime filters** of the algebra.
- ▶ Show relative completeness with respect to that.

For this logic: Dzik, Jarvinen, and Kondo [DJK10, §5].

But this is **non-constructive**, and also not very nice.

Coherent semantics

Replace

- ▶ the poset of worlds by a **distributive lattice** (W, \sqsubseteq)
- ▶ upper sets by (non-prime) **filters**

$F \subseteq W$ is a **filter** just if it is an upper set and

$$1 \in F, \quad x \in F \text{ and } y \in F \text{ imply } x \wedge y \in F$$

$$w \vDash p \stackrel{\text{def}}{\equiv} w \in V(p) \in \text{Filt}(W)$$

$$w \vDash \perp \stackrel{\text{def}}{\equiv} (w = 1)$$

$$w \vDash \varphi \wedge \psi \stackrel{\text{def}}{\equiv} w \vDash \varphi \text{ and } w \vDash \psi$$

$$w \vDash \varphi \vee \psi \stackrel{\text{def}}{\equiv} \exists v_1, v_2. v_1 \wedge v_2 \sqsubseteq w \text{ and } v_1 \vDash \varphi, v_2 \vDash \psi$$

$$w \vDash \varphi \rightarrow \psi \stackrel{\text{def}}{\equiv} \forall v. w \sqsubseteq v \text{ and } v \vDash \varphi \text{ imply } v \vDash \psi$$

This semantics is also sound and complete for intuitionistic logic!

Modalities and dualities

All previous work on modalities carries through, nearly verbatim.

The main duality is now

$$\mathbf{Stable}^{\text{op}} \simeq \mathbf{Coh}$$

between

- ▶ distributive lattices and stable (= \wedge -preserving) maps
- ▶ **coherent frames** (= algebraic cHAs whose compact elements form a sub-lattice) and Scott-continuous, \prod -preserving maps (**not** the usual category from Stone duality)

Then

The coherent semantics and the Heyting algebra semantics are **equi-complete, constructively**.

Categorifying the coherent semantics

Making proofs appear engenders a surprise.

Let \mathcal{C} be a category with finite products and coproducts, which is also ‘**co-distributive**’ **category**, i.e. a category in which

$$a + (c \times d) \cong (a + c) \times (a + d)$$

Then

A two-dimensional coherent semantics is a categorical semantics in a **category of algebras**.

Why? Because ‘filters’ are product-preserving presheaves over \mathcal{C} .

If we regard \mathcal{C} as a Lawvere theory, these are **algebras over \mathcal{C}** .

See Adámek, Rosický, and Vitale [[ARV10](#)].

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